• It is not partially dependent on a candidate key.

Show that every 3NF schema is in 2NF. (*Hint*: Show that every partial dependency is a transitive dependency.)

7.20 Give an example of a relation schema R and a set of dependencies such that R is in BCNF but is not in 4NF.

Exercises

- 7.21 Give a lossless decomposition into BCNF of schema R of Exercise 7.1.
- **7.22** Give a lossless, dependency-preserving decomposition into 3NF of schema *R* of Exercise 7.1.
- 7.23 Explain what is meant by *repetition of information* and *inability to represent information*. Explain why each of these properties may indicate a bad relational-database design.
- **7.24** Why are certain functional dependencies called *trivial* functional dependencies?
- 7.25 Use the definition of functional dependency to argue that each of Armstrong's axioms (reflexivity, augmentation, and transitivity) is sound.
- 7.26 Consider the following proposed rule for functional dependencies: If $\alpha \to \beta$ and $\gamma \to \beta$, then $\alpha \to \gamma$. Prove that this rule is *not* sound by showing a relation r that satisfies $\alpha \to \beta$ and $\gamma \to \beta$, but does not satisfy $\alpha \to \gamma$.
- 7.27 Use Armstrong's axioms to prove the soundness of the decomposition rule.
- 7.28 Using the functional dependencies of Exercise 7.6, compute B^+ .
- **7.29** Show that the following decomposition of the schema *R* of Exercise 7.1 is not a lossless decomposition:

$$(A, B, C)$$

 (C, D, E) .

Hint: Give an example of a relation r(R) such that $\Pi_{A,B,C}(r) \bowtie \Pi_{C,D,E}(r) \neq r$

7.30 Consider the following set F of functional dependencies on the relation schema (A, B, C, D, E, G):

$$A \rightarrow BCD$$

$$BC \rightarrow DE$$

$$B \rightarrow D$$

$$D \rightarrow A$$